

IN THE SPECIFICATION

Please amend the paragraph at page 42, lines 4-15, as follows:

Dth is, for example, $(k+1)(k+l)$. Here k is such a number that the product of $N(1), \dots, N(k)$ is not less than the total number of user IDs (identification information items) ($N(1) \leq N(2) \leq \dots \leq N(M)$), and $[[1]]$ is given by the following formula (1):

$$[1 - \prod_{i=1}^{k+l} 1/N(i)]^S \geq 1 - \epsilon_2 \quad (1)$$

where the range of i that assumes $\prod_{i=1}^{k+l} 1/N(i) \geq 1 - \epsilon_2$ or $i = k+1 \sim (k+l)$, $i = k+1 \sim (k+l)$,

$S = M^{C_{k+1}}$ $S = M^{C_{k+l}}$, and

ϵ_2 represents the rate of error tracing in each user ID of the people responsible for collusive attacks, and satisfies $0 < \epsilon_2 < 1$.

Please amend the paragraph beginning at page 52, line 27, to page 53, line 1, as follows:

Assume that $M = c(k+l)$ $M = c(k+l)$, C is a narrow sense $[M, k, M-k+1]_q$ Reed-Solomon code.

Please amend the paragraph at page 53, lines 2-11, as follows:

If the following formula (2) is satisfied, the Reed-Solomon code C can be made to be a stochastic outer code:

$$[1 - 1/q^k]^S \geq 1 - \epsilon \quad [1 - 1/q^l]^S \geq 1 - \epsilon \quad (2)$$

where $S = M^{C_{k+1}}$ $S = M^{C_{k+l}}$,

$q = N(1) = N(2) = \dots = N(M)$,
and ε represents the rate of error tracing in each user
ID (identification information) of the people responsible for
collusive attacks, and is a real number that satisfies $0 < \varepsilon <$
1.

Please amend the paragraph at page 53, lines 12-17, as follows:

In this case, the above-described tracing algorithm
example as a stochastic method is applicable. In the tracing
algorithm example as a stochastic method, $[[1]]$ is included in
the formula, ~~$D_{th} = k + 1$~~ $D_{th} = k + \ell$, may be given by, for
example, formula (2) instead of formula (1).